



An Integrated Privacy Preserving Attribute Based Access Control Framework

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Cloud Storage

Top Cloud Storage Providers

Google Cloud

Amazon Web Service

Microsoft Azure...

Take MS Azure as an Example

2012, 4 Trillion Objects

2015 Jan, 10 Trillion Objects



Source: <https://www.nasuni.com/infographic-2015-state-of-cloud-storage/>

Cloud Storage

Recent advances have enabled applications that generate/collect *huge amounts of personal data*.

Cloud Storage Providers

Honest-but-Curious

-- *run the programs and algorithms correctly
but gather information related to the stored data and access records.*



“ At year-end 2016, more than **50%** of **Global 1000 companies** will have stored customer-sensitive data in the public cloud ”

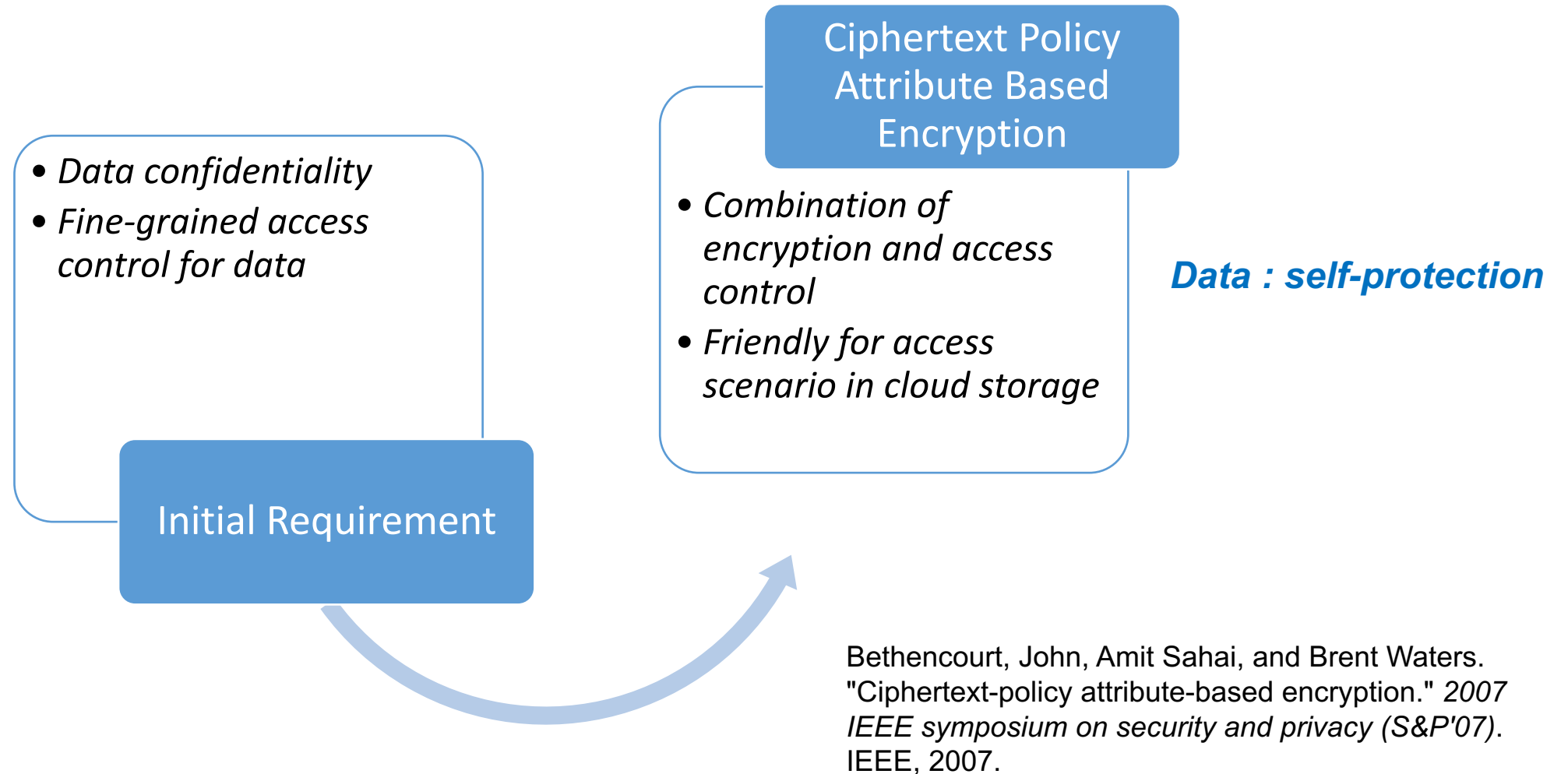
– Gartner

Source: <http://www.gartner.com/newsroom/id/1862714>

Gartner, Inc. is the world's leading information technology research and advisory company.

Security & Privacy Concerns: Personal Data / Sensitive Data

Initial Solution



Scenario

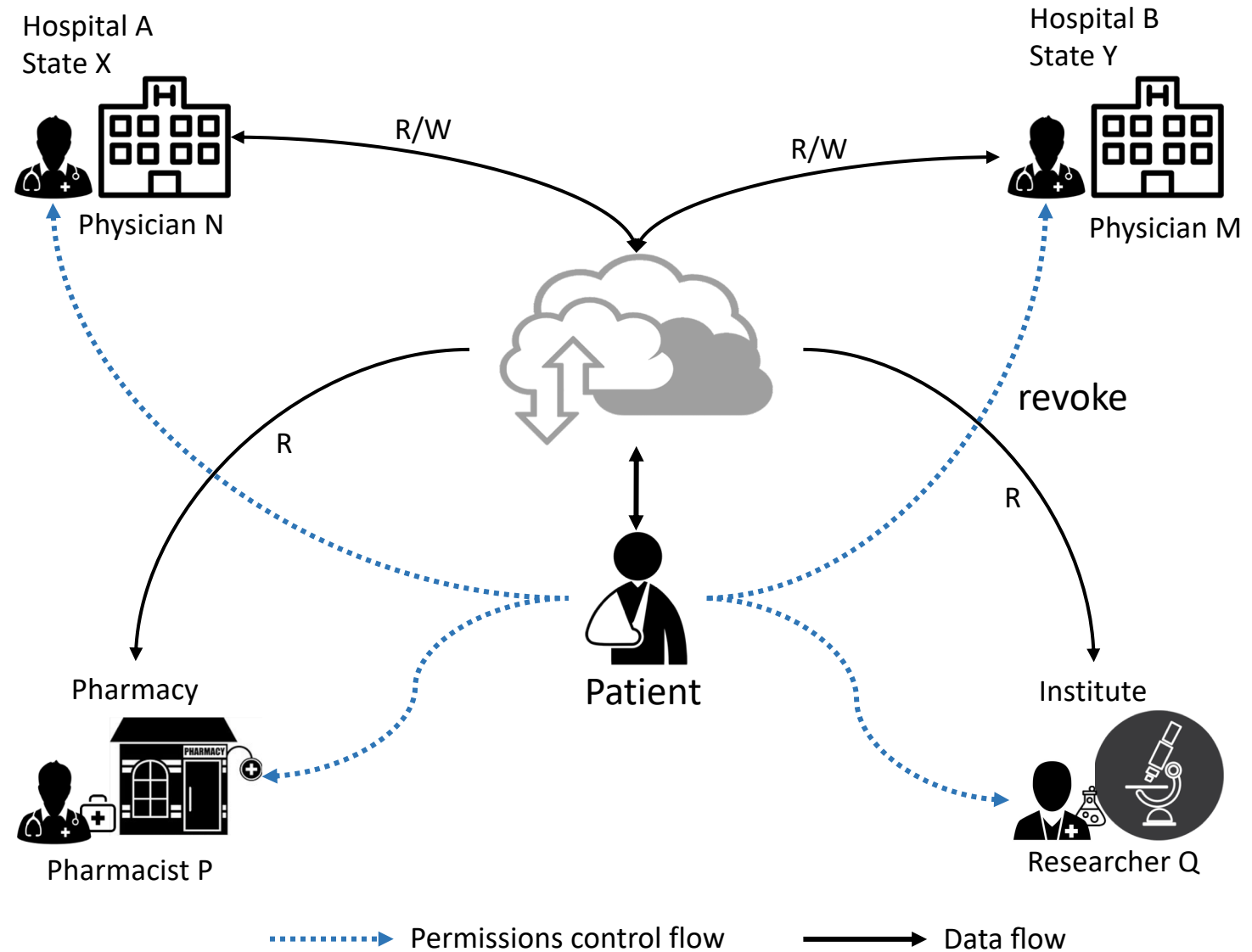
A patient-centric health application

-- that allows a patient/user to store and manage all his Electronic Health Records (EHRs) by storing them in Cloud Storage

Similar scenarios:

- User-centric applications
- Organization-centric applications
- Hospital-centric applications

How to protect user-sensitive data in the public cloud ?



Challenges

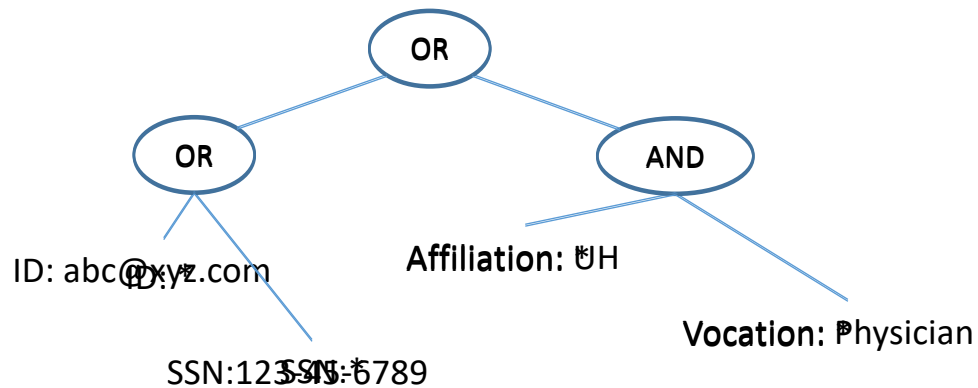
- Challenges of applying CP-ABE to the Scenarios
 - *Support both revocation and privacy-preserving policy*
 - *Limitation of all CP-ABE schemes*
 - Only support read access
 - But don't support write access & policy update
 - *Access patterns leak*
 - Data is protected by encryption, it doesn't matter ?
 - E.g., encrypted data in the cloud, which is often accessed from hospitals, may be identified as EHRs, then link to a specific patient.

Social engineering
Enough time
Powerful machine



The key contributions

- A privacy-preserving revocable CP-ABE scheme (PR-CP-ABE)
 - *Privacy-preserving Access Structure*
 - Linear Secret Sharing Scheme (LSSS)
 - *Supports immediate attribute revocation*



*(ID: abc@xyz.com OR SSN: 123-45-6789) OR
(Affiliation: University Hospital AND Vocation: Physician)*



*(ID: * OR SSN: *) OR (Affiliation: * AND Vocation: *)*

The key contributions

- An extended path oblivious RAM (ePath-ORAM) protocol
 - *Prevents privacy disclosure of access patterns*
 - *Supports data/policy update*
- Security proof of the PR-CP-ABE scheme

Preliminaries: What's CP-ABE

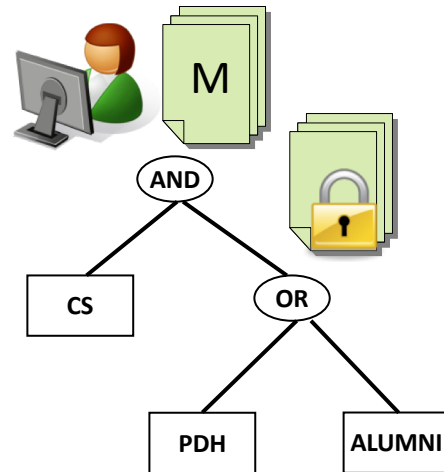
Slide from ESORICS

CP-ABE in detail

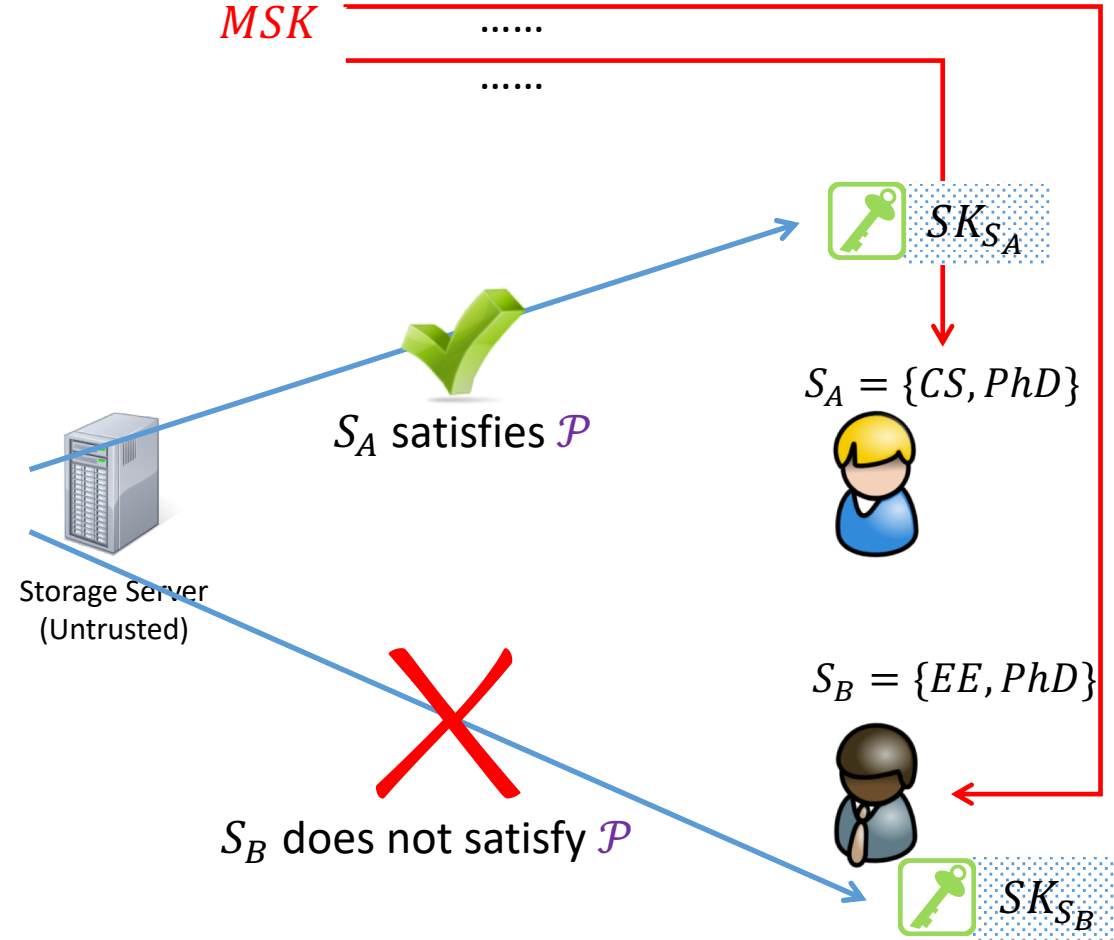
PK PK_{CS}, PK_{EE}, \dots
 $PK_{PhD}, PK_{ALU}, \dots$
 PK_M, PK_F, \dots
 $PK_{1980}, PK_{1981}, \dots$
 \dots

U
 MSK
 Dept.: CS, EE, ...
 Type: PhD Stud., Alumni, ...
 Gender: Male, Female
 Birth Year: 1980, 1981, ...

$$C = Enc(PK, \mathcal{P}, M)$$

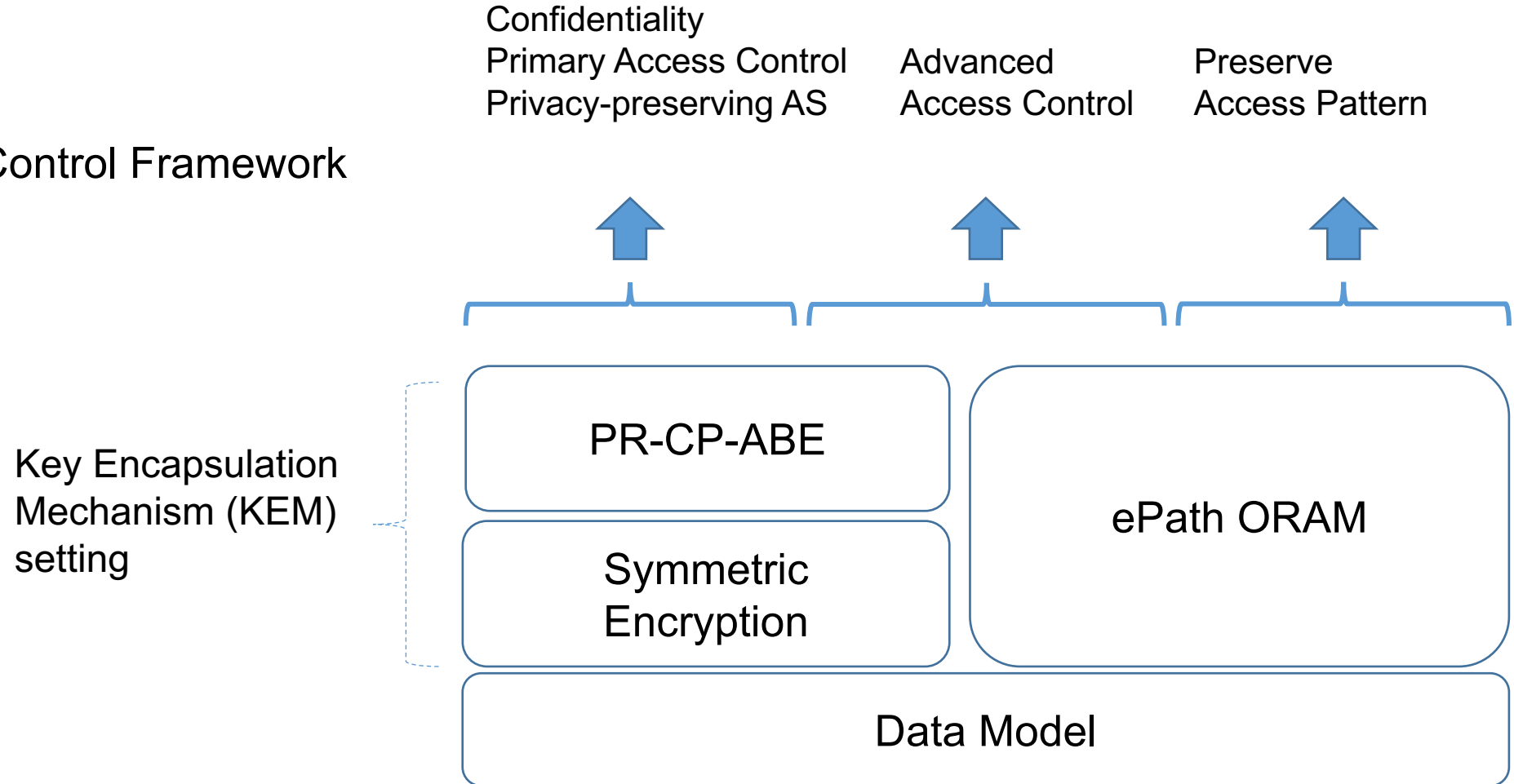


$$\mathcal{P} = CS \text{ AND } (PhD \text{ OR } ALU)$$



Overview

of Access Control Framework



Data Model

$$\mathcal{D} = (id, \mathcal{P}_r, \mathcal{P}_w, \mathcal{P}_o, Enc_{k_\delta}(data)),$$

where

$$\mathcal{P}_r = (\langle A_r, \rho_r \rangle, Enc_\gamma(k_\delta)),$$

$$\mathcal{P}_w = (\langle A_w, \rho_w \rangle, Enc_\gamma(s_w), s_w),$$

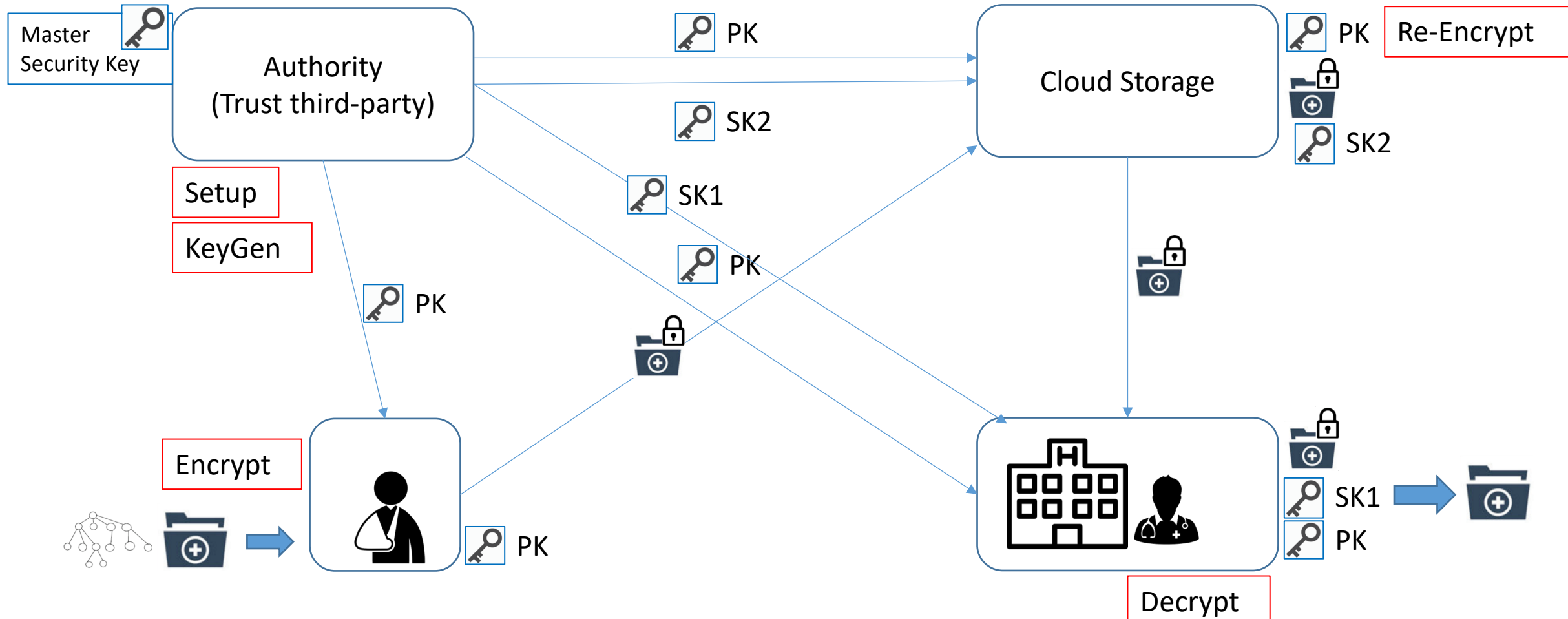
$$\mathcal{P}_o = (\langle A_o, \rho_o \rangle, Enc_\gamma(s_o), s_o).$$

Encrypted data under KEM setting

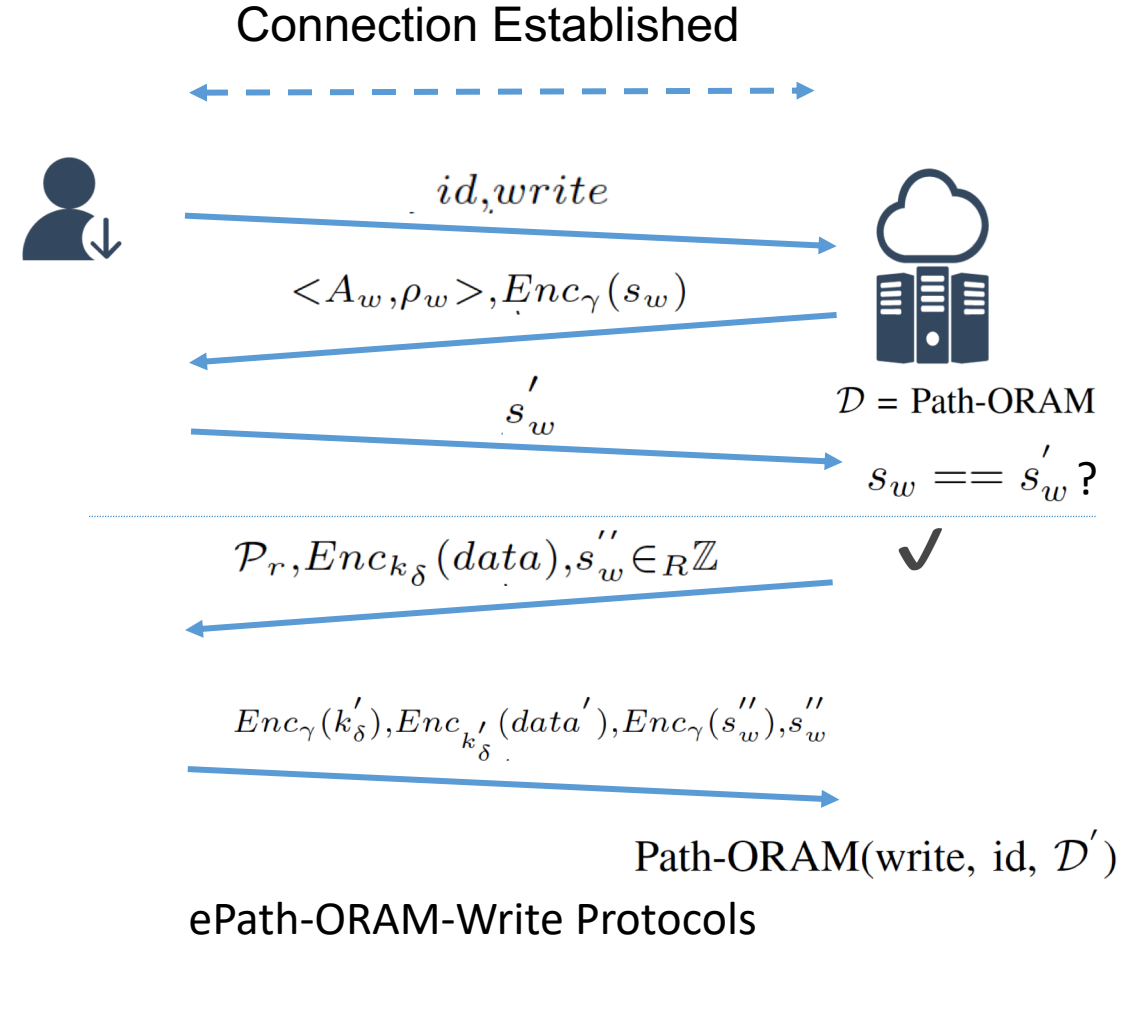
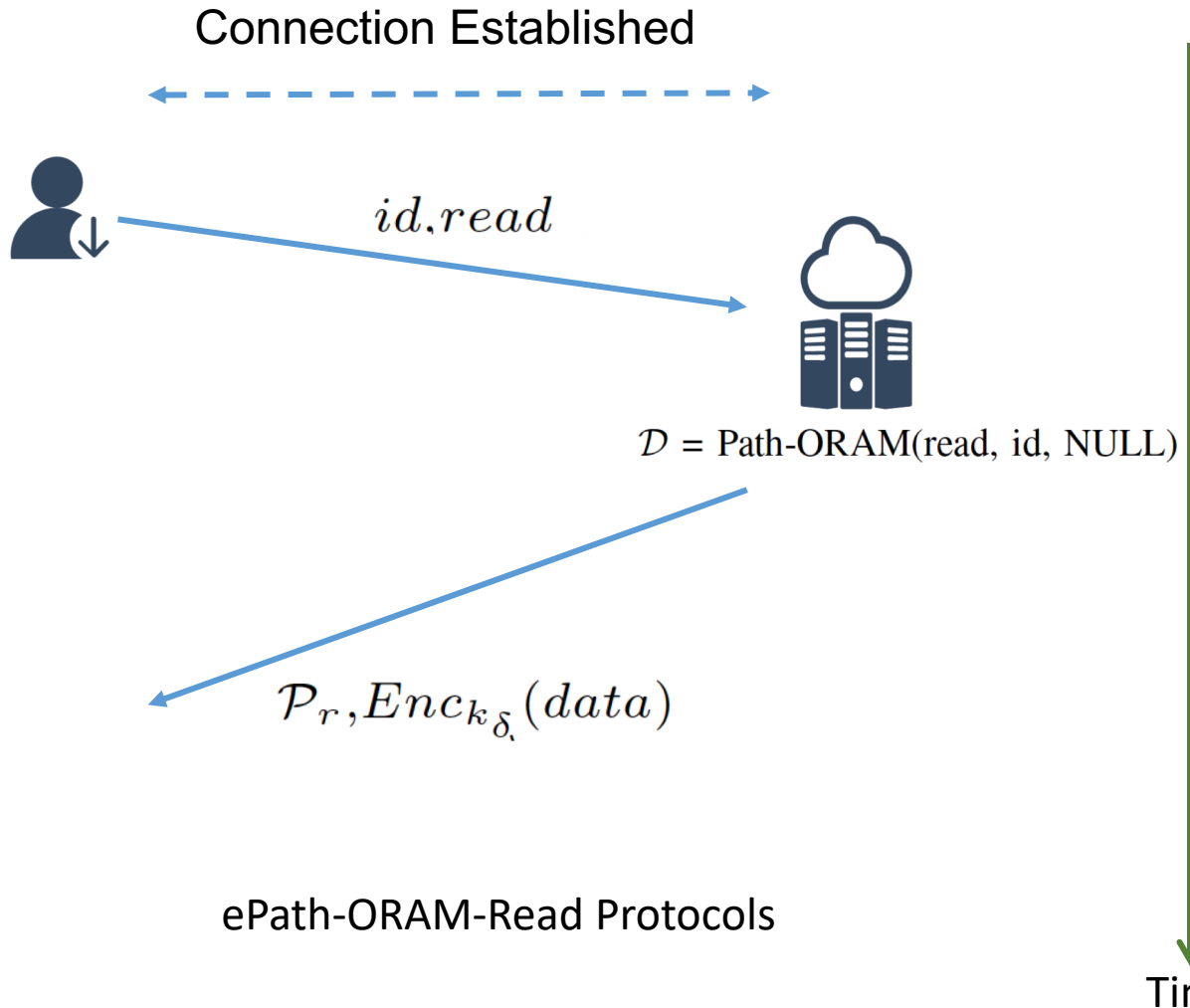
Used to verify a user's write permission
By checking decryption ability on a random
seed

Three access structure (hide value)

Instance of PR-CP-ABE



ePath ORAM Protocols



Analysis

- Tricks behind PR-CP-ABE construction
 - *Composite Order Bilinear Groups*
 - Introduce random elements from a subgroup into algorithms to perturb/hide ciphertext components \leftrightarrow attributes
 - Use the property to eliminate random elements

$$\begin{cases} h_r \in \mathbb{G}_r \\ h_p \in \mathbb{G}_p \end{cases} \longrightarrow e(h_r, h_p) = 1$$

- *Re-encrypt technology*
 - Divide the initial secret element $\alpha = \alpha_1 + \alpha_2$
 - One is corresponding to user, the other is for delegation (CSP)

Analysis

- Forward Security

- *Protects past ciphertext against future compromises of secret keys.*
- *If attribute is revoked*
 - Users can not update the corresponding private key
 - Thus they can not decrypt again

- Backward Security

- *A new user joins in an attribute group that satisfies the policy*
- *Suppose he has a previous ciphertext*
- *Even if he can update private key, he can not decrypt it*
 - Random elements in previous component $D' \leftarrow X \rightarrow$ new user's private key

Analysis

- Key features

Table 1: Comparison of key features

Schemes	Access Structure (AS)	Immediate Revocation	Privacy-preserving AS
[39]	LSSS Matrix	Yes	No
[38]	And-gate	Yes	No
[15]	Tree-based	Yes	No
[20]	LSSS Matrix	No	Yes
[23]	And-gate	No	Yes
Ours	LSSS Matrix	Yes	Yes

Analysis

- Performance

- *As shown in previous experiments.*

- Encryption/decryption → milliseconds level

- Key Application (network communication) → seconds level

Table 2: Comparison of communication cost

Entities	Our scheme	[39]	[38]	[15]
Authority ↔ User	$(2 + n_{a,i}) g $	$(2 + n_{a,i}) g $	$(1 + 2n_{a,i}) g $	$(1 + 2n_{a,i}) g $
Authority ↔ Owner	$(2 + n) g + g_T $	$2 g + g_T $	$(1 + 3n_{a,i}) g + g_T $	$2 g + g_T $
CSP ↔ User	$(4l + 3) g + 2 g_T $	$(2l + 3) g + g_T + l p $	$(3l + 2n_{a,i}) g + g_T $	$(3l + 2n_{a,i}) g + g_T $ $+ (l/2 \times n_u + \log(n_u + 1)) p $
CSP ↔ Owner	$2((2l + 1) g + g_T)$	$(2l + 1) g + g_T $	$3l g + g_T $	$2l g + (l + 1) g_T $

¹ $|p|$, $|g|$ and $|g_T|$ are the elements size in \mathbb{Z}_p , \mathbb{G} and \mathbb{G}_T , respectively.

² $n_{a,i}$ is the number of attributes that the user i possess.

³ l represents the number of attributes associated with the ciphertext.

Our scheme makes a compromise on performance for privacy-preserving policy, compared with [39]
However, our scheme's performance is better than others

Security Proof

- Methodology
 - *Suppose that adversary has advantage to break our scheme*
 - *Adversary's advantage \leftrightarrow break q -parallel BDHE assumption*
 - *However, no-polynomial time algorithm has advantage to break assumption*
 - *Thus no adversary has advantage to break our scheme*

Please find the detail proof in Appendix.

Conclusion

- A novel privacy-preserving attribute-based access control framework
 - *PR-CP-ABE*
 - Privacy-preserving
 - Revocation
 - Security Proof: CPA
 - *ePath-ORAM Protocol*
 - Preserve access pattern
 - Extend PR-CP-ABE to support r/w/o access
 - *Features*
 - User-centric data and policy management
 - Immediate privilege revocation
 - Privacy protection

Q & A

Thanks